

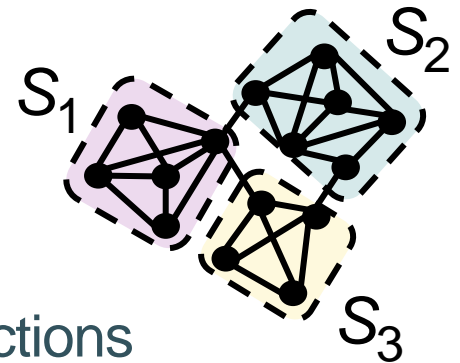
Minimum Average Cost Clustering

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For clustering problems with submodular objective functions, we introduce the minimum average cost criterion

The proposed algorithm

- does not require # of clusters in advance
- computes an optimal # of clusters and an optimal partition in polynomial time
- uses the theory of intersecting submodular functions



Keywords

clustering, [submodular functions](#), combinatorial optimization

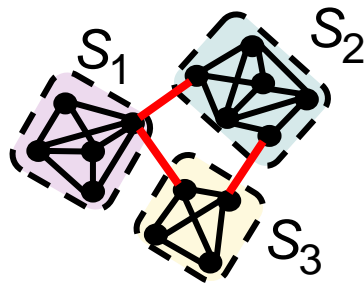
- $V = \{1, \dots, n\}$ is a finite set of data points
- A function f defined on $2^V = \{S: S \subseteq V\}$ is **submodular** if

$$f(S) + f(T) \geq f(S \cup T) + f(S \cap T), \quad \forall S, T \subseteq V$$
 - A generalization of **cut functions**, **entropy functions**, etc.

Clustering problem with submodular objective function

[Narasimhan-Jojic-Bilmes, NIPS 2005]

Given set V , integer $k (\leq n)$, and submodular function f , find a k -partition of V , $\{S_1, \dots, S_k\}$ that minimizes $\sum_{i=1}^k f(S_i)$



optimal 3-clustering

optimal k -clustering

In the case of a network, $\sum_i f(S_i) = 2 \times \#(\text{red edge})$
 (V is a set of nodes, and f is a cut function)

■ Optimal k -clustering problem [Narasimhan *et al.*, NIPS 2005]

$$\begin{aligned} \min \quad & \sum_{i=1}^k f(S_i) \\ \text{s. t.} \quad & \{S_1, \dots, S_k\} \text{ is a } k\text{-partition of } V \end{aligned}$$

○ k (# of clusters) should be computed via some method

☹ NP-hard

■ Minimum Average Cost (MAC) clustering [This work]

$$\begin{aligned} \min \quad & \sum_{S \in \mathcal{P}} f(S) / (|\mathcal{P}| - \beta) \\ \text{s. t.} \quad & \mathcal{P} \text{ is a partition of } V \\ & |\mathcal{P}| > \beta \end{aligned}$$

averaged objective function

If $\beta = 1$, then the objective function is a natural average cost of \mathcal{P}

where $0 \leq \beta < n$

β -MAC clustering

○ $k = |\mathcal{P}|$ and a partition \mathcal{P} are determined at the same time

😊 Solvable in poly time

😊 competitive with other methods

If β is small, \mathcal{P} is coarse. If β is big, \mathcal{P} is fine.

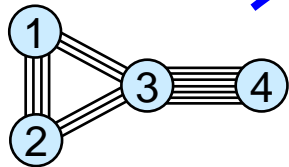
Theory of intersecting submodular functions

⇒ **Theorem** [This work]. There is an algorithm that computes all the β -MAC clusterings in **polynomial time** in total

Observation. Suppose that a partition \mathcal{P} is a β -MAC clustering for some β , and let $k = |\mathcal{P}|$. Then, \mathcal{P} is a k -optimal clustering

⇒ The information about MAC clusterings gives a portion of the information about optimal k -clusterings

(Remember that an optimal k -clustering problem is NP-hard)

Example

Compute
all β -MAC
clusterings

$$0 \leq \beta < 1$$

$$\{\{1, 2, 3, 4\}\}$$

opt 1-clustering

$$1 \leq \beta < 11/7$$

$$\{\{1\}, \{2\}, \{3, 4\}\}$$

opt 3-clustering

$$11/7 \leq \beta < 4$$

$$\{\{1\}, \{2\}, \{3\}, \{4\}\}$$

opt 4-clustering

In this case, our algorithm computes optimal k -clusterings for $k = 1, 3, \& 4$

For more information,
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