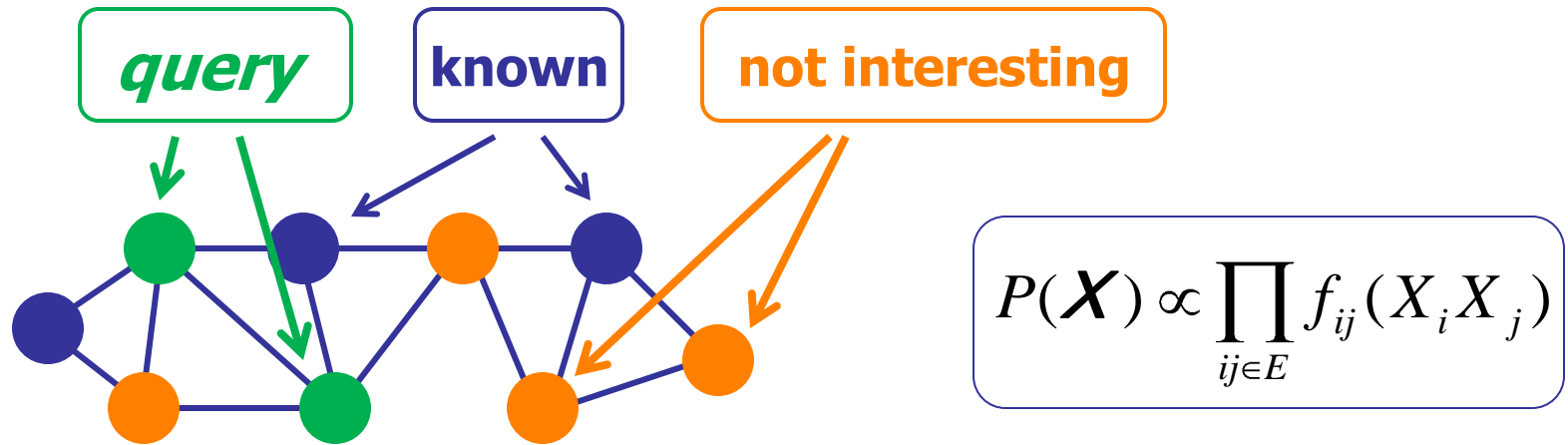


# Focused Belief Propagation for Query-Specific Inference

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# Query-Specific inference problem



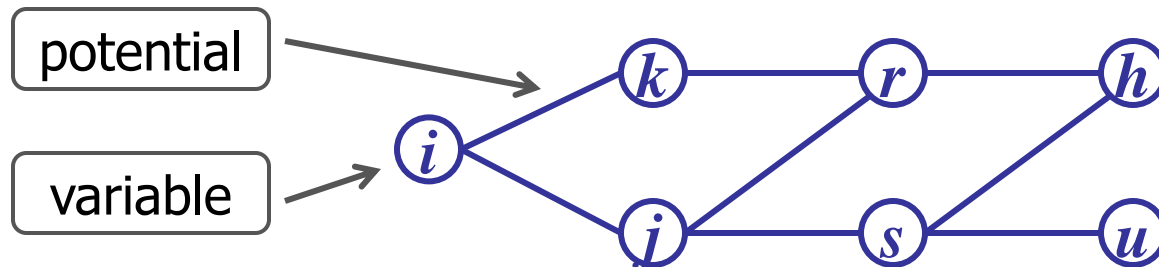
Using information about **the query**  
to speed up convergence of **belief propagation**  
*for the query marginals*

# Graphical models

- **Talk focus: *pairwise Markov random fields***
- **Paper: *arbitrary factor graphs*** (more general)
- Probability is a product of **local** potentials

$$P(\mathbf{X}) \propto \prod_{ij \in E} f_{ij}(X_i X_j)$$

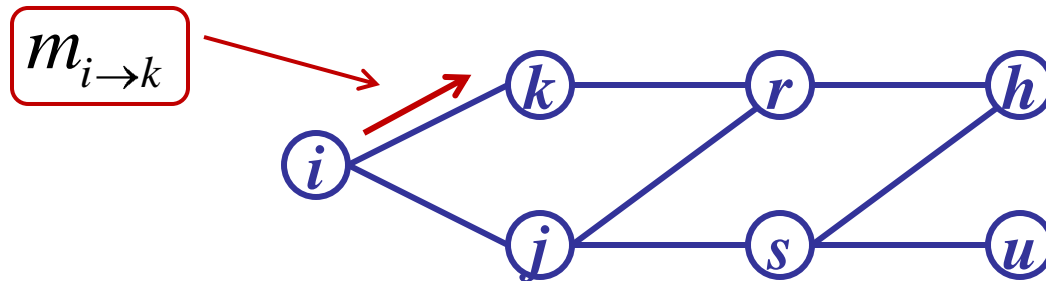
- Graphical structure



- ☺ Compact representation
- ☹ Intractable inference
  - ☺ Approximate inference often works well in practice

# (loopy) Belief Propagation

- Passing messages along edges



- Variable belief:

$$\tilde{P}^{(t)}(x_i) \propto \prod_{ij \in E} m_{j \rightarrow i}^{(t)}(x_i)$$

- Update rule:

$$m_{j \rightarrow i}^{(t+1)}(x_i) = \sum_{x_j} f_{ij}(x_i, x_j) \prod_{kj \in E, k \neq i} m_{k \rightarrow j}^{(t)}(x_j)$$

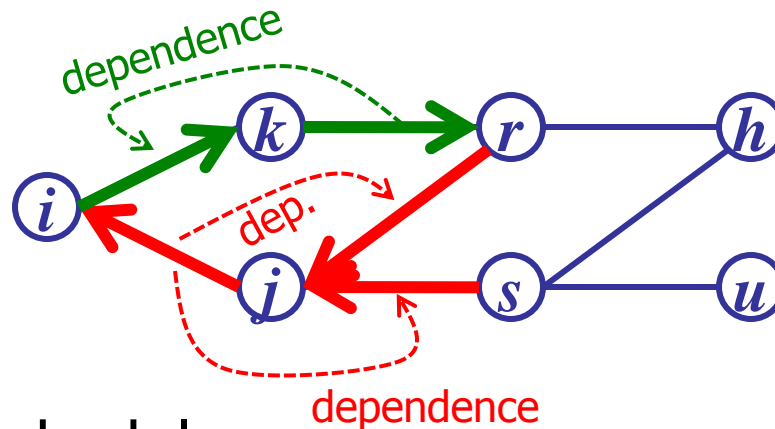
- Result: all single-variable beliefs

# (loopy) Belief Propagation

- Update rule:

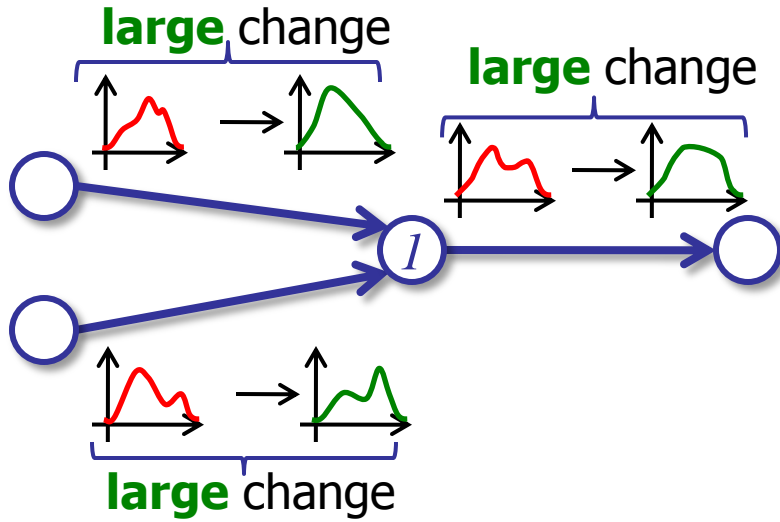
$$m_{j \rightarrow i}^{(t+1)}(x_i) = \sum_{x_j} f_{ij}(x_i, x_j) \prod_{k \in E, k \neq i} m_{k \rightarrow j}^{(t)}(x_j)$$

- Message dependencies are local:

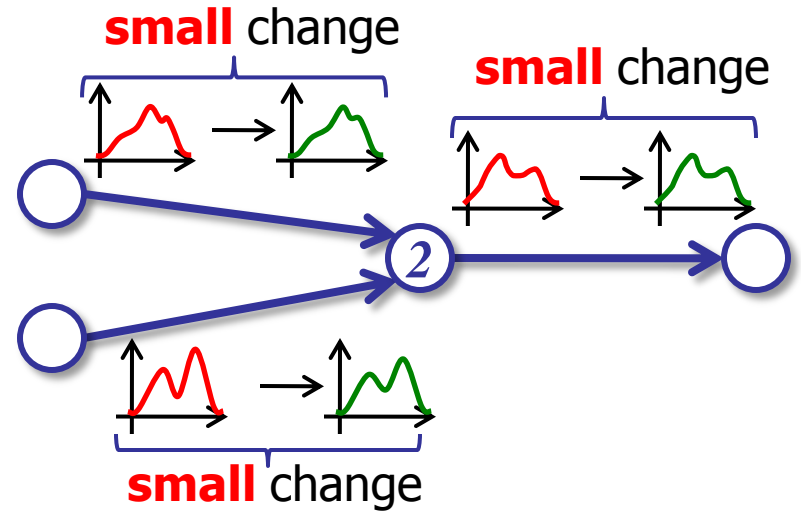


- Round-robin schedule
  - Fix message order
  - Apply updates in that order until convergence

# Dynamic update prioritization



informative update



wasted computation

- **Fixed** update **sequence** is not the best option
- Dynamic **update scheduling** can speed up convergence
  - Tree-Reweighted BP [Wainwright et. al., AISTATS 2003]
  - Residual BP [Elidan et. al. UAI 2006]
- Residual BP → apply the **largest change** first

# Residual BP [Elidan et. al., UAI 2006]

- Update rule:

$$\text{new} \leftarrow m_{j \rightarrow i}^{(NEW)}(x_i) = \sum_{x_j} f_{ij}(x_i x_j) \prod_{k \in E, k \neq i} m_{k \rightarrow j}^{(OLD)}(x_j) \rightarrow \text{old}$$

- Pick edge with ***largest residual***

$$\max \left\| m_{j \rightarrow i}^{(NEW)} - m_{j \rightarrow i}^{(OLD)} \right\|$$

- Update  $m_{j \rightarrow i}^{(OLD)} \leftarrow m_{j \rightarrow i}^{(NEW)}$

More effort on the difficult parts of the model 😊

But no query ☹️

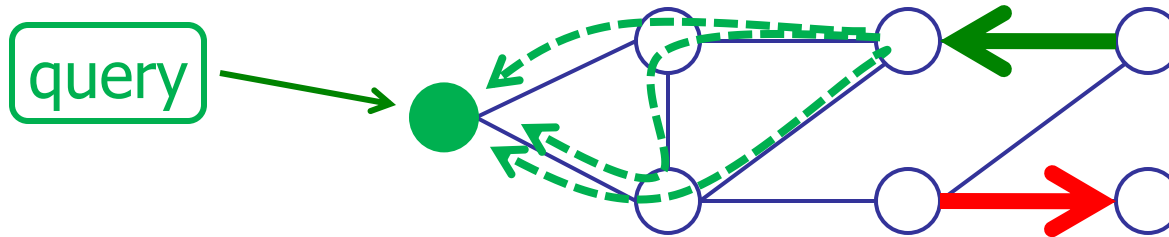
# Our contributions

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- Using ***weighted*** residuals to prioritize updates
- Define message weights reflecting the ***importance*** of the message ***to the query***
- Computing importance weights efficiently
- Experiments: faster convergence on large relational models



# Why edge importance weights?



**residual** < **residual**  
which to update??

- Residual BP updates →
- ***no influence*** on the query
- wasted computation

- want to update ←
- ***influence*** on the query  
***in the future***

Our work → max approx. ***eventual effect on  $P(query)$***

Residual BP → max ***immediate residual reduction***

# Query-Specific BP

- Update rule:

$$\text{new} \leftarrow m_{j \rightarrow i}^{(NEW)}(x_i) = \sum_{x_j} f_{ij}(x_i x_j) \prod_{kj \in E, k \neq i} m_{k \rightarrow j}^{(OLD)}(x_j) \rightarrow \text{old}$$

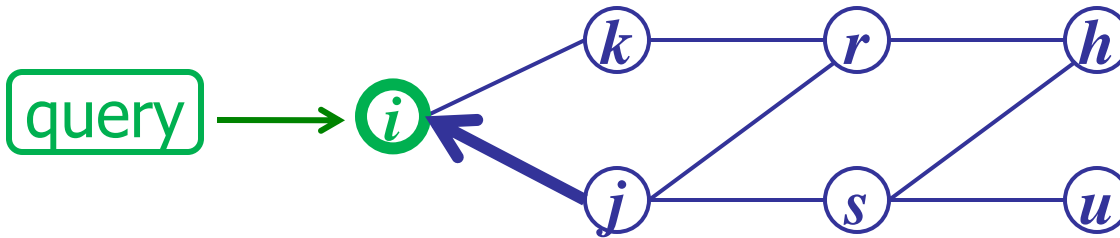
- Pick edge with  $\max \left\| m_{j \rightarrow i}^{(NEW)} - m_{j \rightarrow i}^{(OLD)} \right\| \rightarrow \text{edge importance}$
- Update  $m_{j \rightarrow i}^{(OLD)} \leftarrow m_{j \rightarrow i}^{(NEW)}$  the only change!

Rest of the talk:  
***defining and computing edge importance***

# Edge importance base case

- Pick edge with  $\max \left\| m_{j \rightarrow i}^{(NEW)} - m_{j \rightarrow i}^{(OLD)} \right\| \times A_{j \rightarrow i}$

approximate eventual update effect on  $P(\mathbf{Q})$



Base case: edge **directly connected** to the query

$$A_{j \rightarrow i} = ??$$

change in query belief

change in message

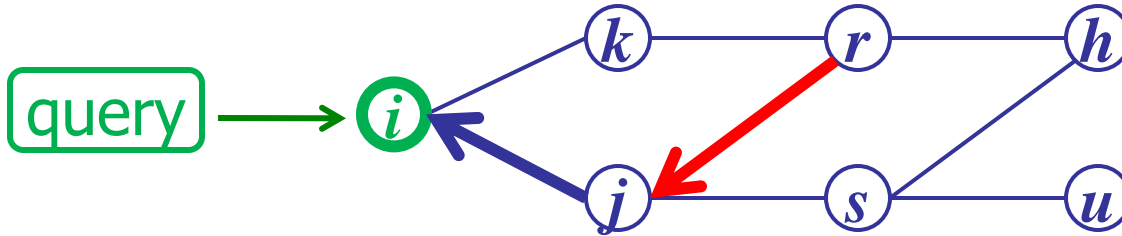
$$\left\| P^{(NEW)}(\mathbf{Q}) - P^{(OLD)}(\mathbf{Q}) \right\| \leq \left\| m_{j \rightarrow i}^{(NEW)} - m_{j \rightarrow i}^{(OLD)} \right\| \times \mathbf{1}$$

$$\left\| \Delta P(\mathbf{Q}) \right\|$$

tight bound

$$\left\| \Delta m_{j \rightarrow i} \right\|$$

# Edge importance one step away

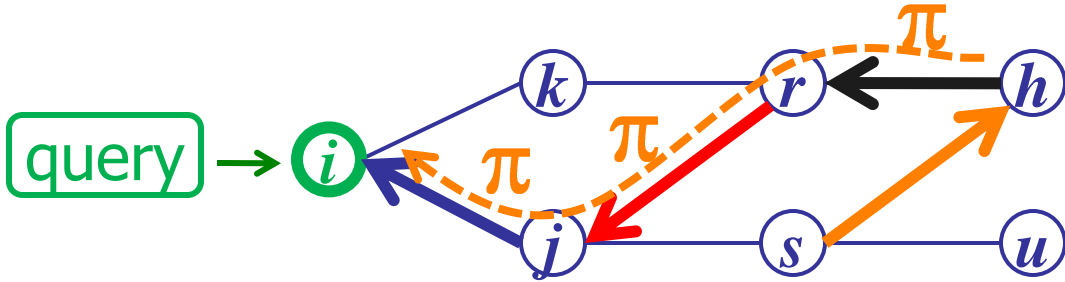


Edge **one step away** from the query:  $\mathbf{A}_{r \rightarrow j} = ??$

$$\underbrace{\|\Delta P(\mathbf{Q})\|}_{\text{change in query belief}} \leq \|\Delta \mathbf{m}_{j \rightarrow i}\| \leq \underbrace{\|\Delta \mathbf{m}_{r \rightarrow j}\|}_{\text{change in message}} \times \sup_{\text{over values of all other messages}} \underbrace{\left\| \frac{\partial \mathbf{m}_{i \rightarrow i}}{\partial \mathbf{m}_{r \rightarrow j}} \right\|}_{\text{message importance}}$$

can compute **in closed form**  
looking at only  $\mathbf{f}_{ji}$  [Mooij, Kappen; 2007]

# Edge importance general case



Base case:  $A_{j \rightarrow i} = 1$

One step away:

$$A_{r \rightarrow j} = \sup \left\| \frac{\partial m_{j \rightarrow i}}{\partial m_{r \rightarrow j}} \right\|$$

$$\|\Delta P(\mathbf{Q})\| \leq \|\Delta m_{s \rightarrow h}\| \times \sup \left\| \frac{\partial P(\mathbf{Q})}{\partial m_{s \rightarrow h}} \right\|$$

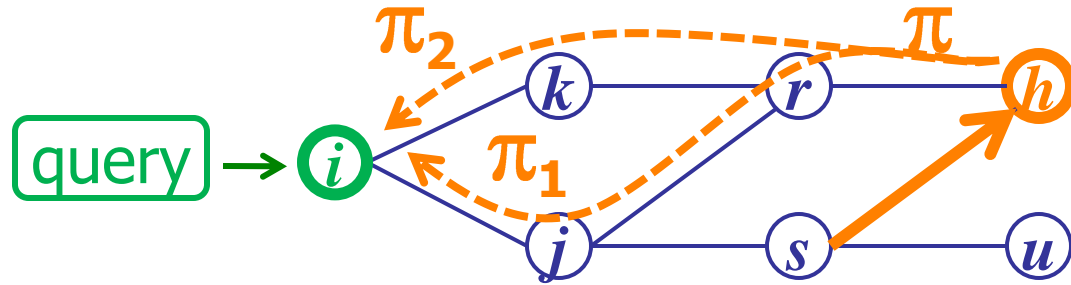
Generalization?

- ☹ *expensive* to compute
- ☹ bound may be *infinite*

$$\sup \left\| \frac{\partial P(\mathbf{Q})}{\partial m_{s \rightarrow h}} \right\|_{\pi} \leq \underbrace{\sup \left\| \frac{\partial m_{h \rightarrow r}}{\partial m_{s \rightarrow h}} \right\| \times \sup \left\| \frac{\partial m_{r \rightarrow j}}{\partial m_{h \rightarrow r}} \right\| \times \sup \left\| \frac{\partial m_{j \rightarrow i}}{\partial m_{r \rightarrow j}} \right\|}_{\text{sensitivity}(\pi)}$$

**sensitivity**( $\pi$ ): max impact *along the path*  $\pi$

# Edge importance general case



$$\sup \left\| \frac{\partial P(\mathbf{Q})}{\partial \mathbf{m}_{s \rightarrow h}} \right\|_{\pi} \leq \underbrace{\sup \left\| \frac{\partial \mathbf{m}_{h \rightarrow r}}{\partial \mathbf{m}_{s \rightarrow h}} \right\| \times \sup \left\| \frac{\partial \mathbf{m}_{r \rightarrow j}}{\partial \mathbf{m}_{h \rightarrow r}} \right\| \times \sup \left\| \frac{\partial \mathbf{m}_{j \rightarrow i}}{\partial \mathbf{m}_{r \rightarrow j}} \right\|}_{\text{sensitivity}(\pi)}$$

**sensitivity**( $\pi$ ): max impact *along the path*  $\pi$

$$\mathbf{A}_{s \rightarrow h} = \mathbf{max} \text{ all paths } \pi \text{ from } \textcircled{h} \text{ to query } \text{sensitivity}(\pi)$$

There are **a lot** of paths in a graph,  
trying out every one is intractable ☹️

# Efficient edge importance computation

$$A = \max_{\text{all paths } \pi \text{ from } \leftarrow \text{ to query}} \text{sensitivity}(\pi)$$

There are **a lot** of paths in a graph,  
trying out every one is intractable ☹️

*always decreases* as the path grows

$\text{sensitivity}(h \rightarrow r \rightarrow j \rightarrow i) =$

**Dijkstra's** (shortest paths) **alg.**  
will *efficiently* find max-sensitivity paths  
for every edge 😊

*decomposes* into  
individual edge contributions

# Query-Specific BP

- Run Dijkstra's alg **starting at query** to get edge weights

$$A_{j \rightarrow i} = \max_{\text{all paths } \pi \text{ from } i \text{ to query}} \text{sensitivity}(\pi)$$

More effort on the difficult **and relevant** parts of the model

$$\max \left( \|m_{j \rightarrow i}\| - \|m_{j \rightarrow i}\| \right) \times A_{j \rightarrow i}$$

- Takes into account not only ***graphical structure***, but also ***strength of dependencies***



- Using *weighted* residuals to prioritize updates
- Define message weights reflecting the *importance* of the message *to the query*
- Computing importance weights efficiently
  - As an initialization step before residual BP
- Restore *anytime behavior* of residual BP
- Experiments: faster convergence on large relational models

# Big picture

Before



all  
marginals



anytime  
property

all  
marginals



all  
marginals



After



long  
initialization



Dijkstra



BP



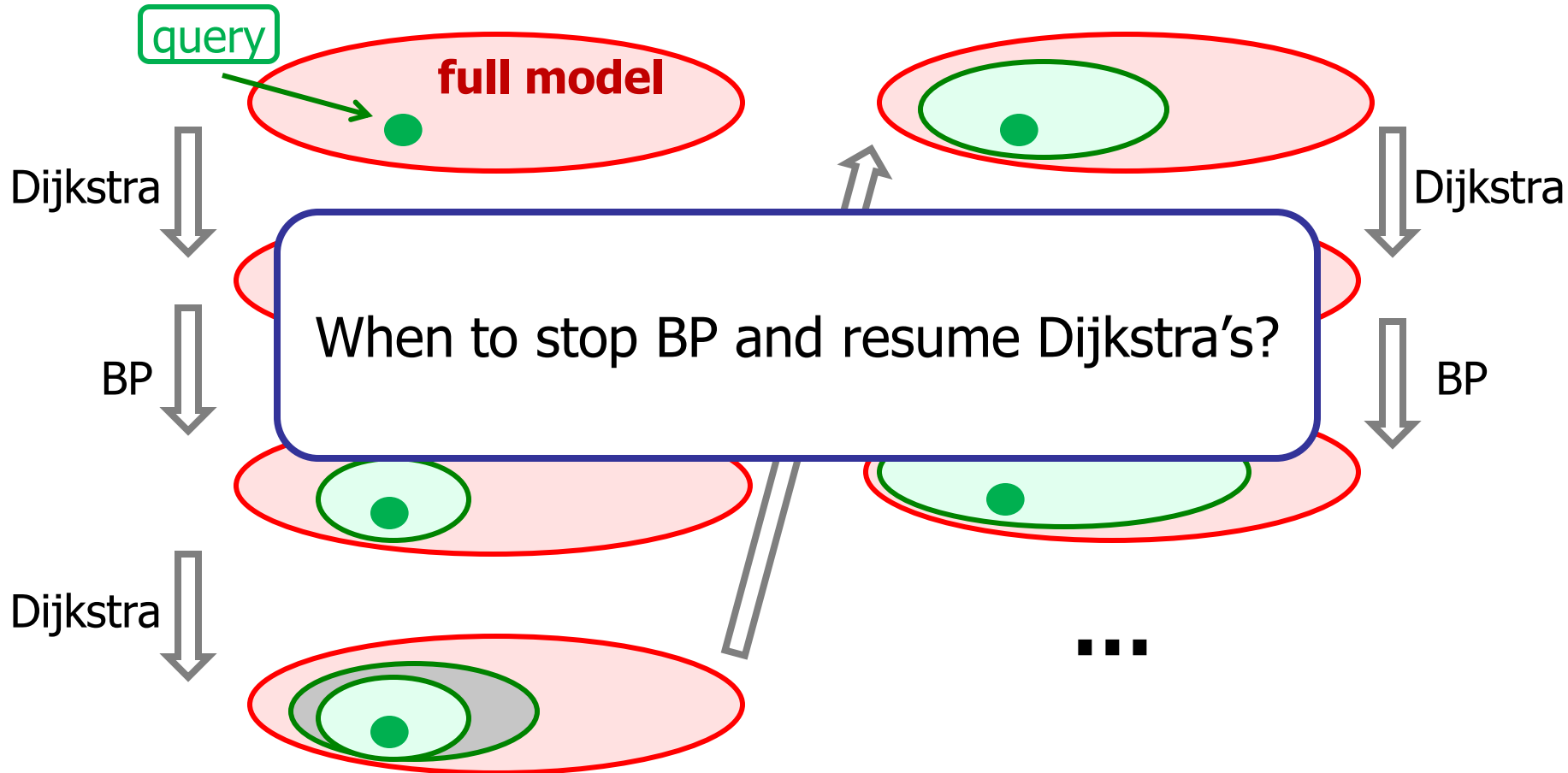
query  
marginals



This is broken!

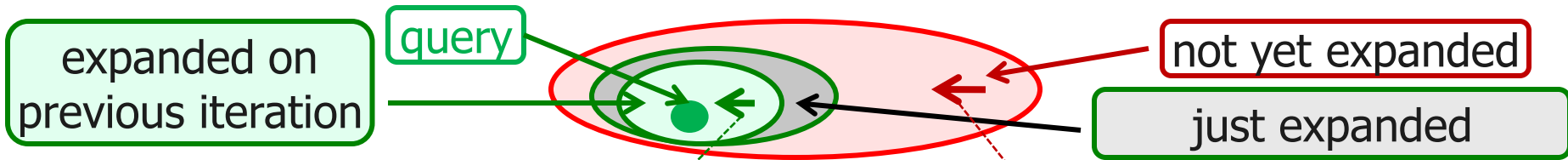
# Interleaving BP and Dijkstra's

- Dijkstra's **expands** the **highest weight** edges **first**
  - Can pause it at any time and get the **most relevant submodel**



# Interleaving

- Dijkstra's expands the **highest weight edges first**



$$\min_{\leftarrow \in \text{expanded edges}} A_{\leftarrow} \geq A_{\leftarrow}$$

suppose  $M \geq \max_{j \rightarrow i \in \text{ALL EDGES}} \|m_{j \rightarrow i}^{(NEW)} - m_{j \rightarrow i}^{(OLD)}\|$

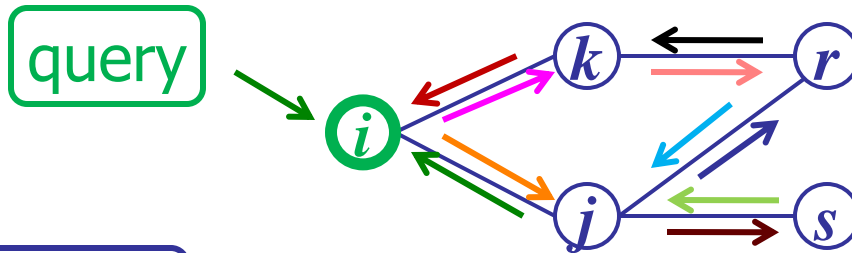
actual priority of  $\leftarrow$

upper bound on  $\leftarrow$  priority

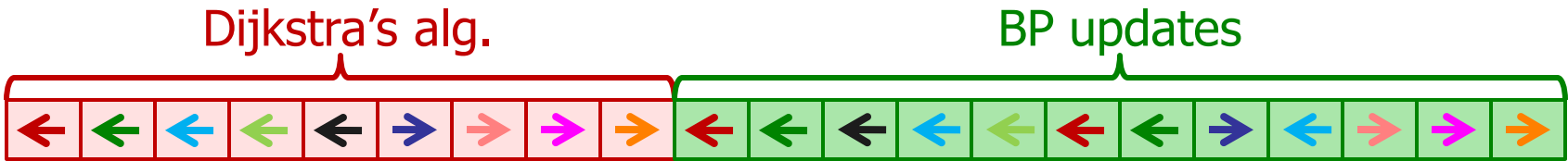
$$\max_{j \rightarrow i \in \text{EXPANDED}} \|m_{j \rightarrow i}^{(NEW)} - m_{j \rightarrow i}^{(OLD)}\| \times A_{j \rightarrow i} \quad \mathbf{M} \times \min_{\leftarrow \in \text{expanded}} A_{\leftarrow}$$

no need to expand further at this point

# Any-Time query-specific BP



Query-specific BP:



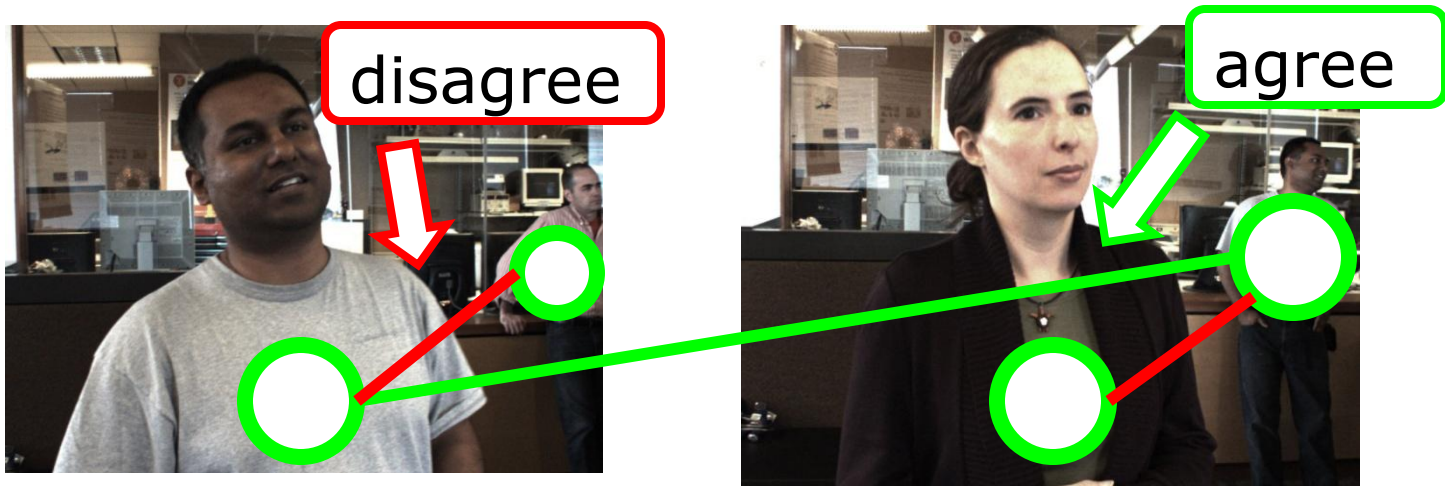
same BP update sequence!

Anytime QSBP:

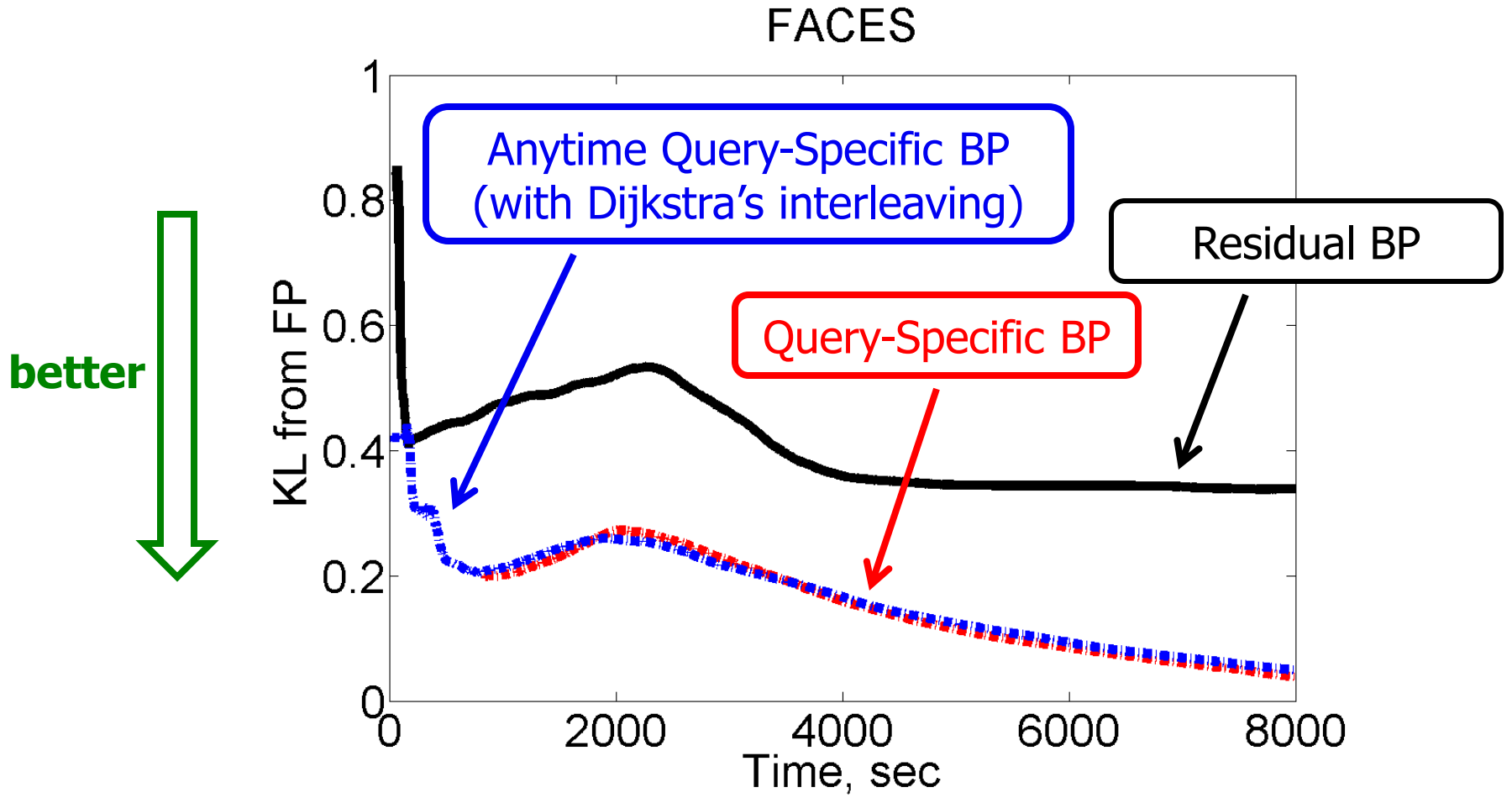


# Experiments – the setup

- Relational model: semi-supervised people recognition in collection of images [with Denver Dash and Matthai Philipose @ Intel]
  - One variable per person
  - Agreement potentials for people who look alike
  - Disagreement potentials for people in the same image
  - 2K variables, 900K factors
- Error measure: KL from the fixed point of BP



# Experiments – convergence speed



# Conclusions

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- Prioritize updates by ***weighted*** residuals
  - Takes ***query info*** into account
- Importance weights depend on both the ***graphical structure*** and ***strength of local dependencies***
- ***Efficient computation*** of importance weights
- Much ***faster convergence*** on large relational models

Thank you!