

Effective Heuristics and Belief Tracking for Planning with Incomplete Information

Alexandre Albore¹ Miquel Ramírez¹ Hector Geffner²

DTIC
Universitat Pompeu Fabra

ICREA & DTIC
Universitat Pompeu Fabra

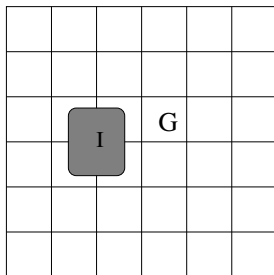
Freiburg, 15/06/2011



Incomplete Information: Conformant Planning

Task

Robot must move from **uncertain** I into G with **certainty**.



- Similar to classical planning except for **uncertain** I
- **Plans**, however, quite **different**: best **conformant** plan **must move the robot to a corner first (localization)**

Conformant planning is **special case** of contingent planning

Ideas provide the **basis** for

- Planning with Sensing (Brafman & Hoffmann, 2005)
- Derivation of Finite–State Controllers (Bonet, Palacios, Geffner, 2009)

Obtaining conformant planners that **scale up** well is critical.

Conformant Planning: Belief State Formulation

- **Belief state:** set of **possible** states
- Actions map belief state b into belief state b_a

$$b_a = \{s' \mid s' \in F(a, s) \ \& \ s \in b\}$$

- Conformant planning is **path-finding** in **belief space**

Challenges: # of belief states is **doubly exponential** in # of vars.

- **Effective representation** of belief states b
- **Effective heuristic** $h(b)$ for estimating cost to G in belief space

Recent alternative: translation into classical planning (Palacios & Geffner 2007).

Basic Translation: Move to the “knowledge level”

conformant problem $P = \langle F, O, I, G \rangle$

- 1 F fluents in P
- 2 O actions with effects $C \rightarrow L$
- 3 I initial situation, CNF **clauses** over F -literals
- 4 G goal situation, **conjunction** of F -literals

classical problem $K_0(P) = \langle F', O', I', G' \rangle$

- 1 $F' = \{KL, K\neg L \mid L \in F\}$
- 2 $O' = O$ **but** preconditions L replaced by KL , effects $C \rightarrow L$ replaced by rules
 - **Support** $KC \rightarrow KL$
 - **Cancellation** $\neg K\neg C \rightarrow \neg K\neg L$
- 3 $I' = \{KL \mid \text{clause } L \in I\}$
- 4 $G' = \{KL \mid L \in G\}$

$K_0(P)$ is **sound** but **incomplete**:

- **Soundness**: every classical plan for $K_0(P)$ is a conformant plan for P
- **Completeness**: all plans for P are classical plans for $K_0(P)$.

Key Elements in General Translation $K_{T,M}(P)$

1 A set T of **tags** t

- consistent sets of **assumptions** (literals) about initial situation I :

$$I \not\models \neg t$$

2 A set M of **merges** m

- **valid subsets of tags** (DNF)

$$I \models \bigvee_{t \in m} t$$

3 Tagged literals KL/t meaning that L true **if** t **initially** true

General Translation: $K_{T,M}(P)$

conformant problem $P = \langle F, O, I, G \rangle$

...

classical problem $K_{T,M}(P) = \langle F', O', I', G' \rangle$

- $F' = \{KL/t, K\neg L/t \mid L \in F \ \& \ t \in T\}$
- $O' = O$ but preconditions L replaced by KL , effects $C \rightarrow L$ replaced by rules
 - **Support** $KC/t \rightarrow KL/t$
 - **Cancellation** $\neg K\neg C/t \rightarrow \neg K\neg L/t$
 - Plus **merge actions**

$$\bigwedge_{t \in M, m \in M} KL/t \rightarrow KL$$

- $I' = \{KL/t \mid \text{if } I \models t \supset L\}$
- $G' = \{KL \mid L \in G\}$

Compiling Uncertainty Away: Properties

- Translation $K_{T,M}(P)$ **always sound**, for suitable choice of sets of **tags** and **merges**, it is **complete**
- **Conformant width** is roughly the **max # of relevant uncertain** variables that interact in P
- $K_i(P)$ is **polynomial instance** of $K_{T,M}(P)$ that is **complete** for problems with conformant width **bounded** by i
- Most benchmarks have **bounded** width and **equal** to 1!
- $K_i(P)$ with $i = 1$, is basis for conformant planner $T0$ (Palacios & Geffner, 2009)

Shortcomings of the Translation–based Approach

- 1 For problems with high width, **complete translation** unfeasible
- 2 **Incomplete** yet **tractable** translations may
 - render a solvable problem **unsolvable**
 - result in **infinite** heuristic values for solvable beliefs
- 3 Relevant information like **cardinality of beliefs**, seems to get **lost in translation**

- 1 **New translation** $K_S^i(P)$
 - Exponential in i , always **complete**, not always **sound**
 - $K_S^i(P)$ **sound** for problems with conformant width $\leq i$
- 2 **New planner** $T1$ based on $K_S^1(P)$ **improves** upon $T0$ planner based on $K_1(P)$
 - **Belief space** planner
 - Two heuristics
 - **Reachability** heuristic h_C derived from $K_S^1(P)$
 - **Certainty** heuristic h_K

Outline for the rest of the talk

- 1 **Translation** $K_S^i(P)$
- 2 **Planner** $T1$
- 3 **Experimental Results**

Idea for the $K_S^i(P)$ translation

$K_{S_0}(P)$ is $K_{T,M}(P)$ with $T = S_0$.

- $K_{S_0}(P)$ sound and complete, but **exponential** on $|F|$

Define K_S to be like K_{S_0} , **but** $T = S \subseteq S_0$

- S **set of samples** of S_0
- $K_S(P)$ is **complete** but **not necessarily sound**.

$K_S^i(P)$ like $K_S(P)$ but with a **specific** sample set S

- **Sound** when width of problem P **bounded** by i
- $|S|$ **exponential** on i

Bases for Conformant Problems

Definition

A **set of states** S , $S \subseteq S_0$ is a **basis for problem** P , iff any conformant plan that conforms with S also conforms with S_0

Theorem (Palacios & Geffner, JAIR 2009)

If problem P has width i , then there exists a basis S for P of size **exponential in i** .

The $K_S^i(P)$ translation

$K_S^i(P)$ is $K_S(P)$ with **sample set** $S \subseteq S_0$ s.t. S guaranteed to be a basis for P if $\text{width}(P) \leq i$

$K_S^i(P)$ is **always** complete and **sound** if $\text{width}(P) \leq i$,

Computation of sample sets S **exponential** in i , provided that I compiled into d -DNNF (Darwiche, 2002)

See paper for details

Belief space planner T1 **implicitly** represents **beliefs**.

Search **node** $n = \langle \pi, S_n, R \rangle$

- π is the **plan prefix** to reach n from **root node** $\langle \emptyset, S^1, R_0 \rangle$
- S_n is sample set S^1 progressed through π
- R set of **known** literals

SAT solver used to check literals true in R

Two heuristics:

- $h_C(n) = h(K_S(P))$, where $S = S_n$
- $h_K(n)$... in next slide

Given node $n = \langle \pi, S_n, R \rangle$, $h_K(n)$ defined as:

of literals L in *one of invariants* overlapping G s.t. $\neg L \notin R$

Related to

- **Landmark** heuristic (Richter, Helmert & Westphal, 2008)
- **Belief cardinality** heuristic (Bertoli & Cimatti, 2002)

T1 planner: Search Engine

- Multi-queue best first search algorithm (Helmert, 2004)
- 3 open lists Q1, Q2, Q3:
 - **Q1**: nodes for helpful actions or that decrease certainty heuristic h_K , ordered with h_C
 - **Q2**: nodes for helpful actions or that decrease certainty heuristic h_K , ordered with h_K
 - **Q3**: nodes for non-helpful actions, ordered with h_C
- We alternate expansion from Q1 and Q2, $1/10$ of the expansions are from Q3.

Experimental Evaluation

		<i>T0</i>			<i>DNF</i>			<i>T1</i>		
<i>Domain</i>	<i>l</i>	<i>S</i>	<i>avg T</i>	<i>avg L</i>	<i>S</i>	<i>avg T</i>	<i>avg L</i>	<i>S</i>	<i>avg T</i>	<i>avg L</i>
bomb	8	8	1.7	111	8	4.0	93	7	0.6	93
coins	8	7	0.1	83	7	0.8	81	8	0.7	142
comm	8	8	0.2	155	7	255.3	170	8	48.4	162
corners-cube	10	8	4.4	303	10	1.4	184	10	95.8	432
cube	6	3	19.6	216	6	959.7	1346	6	6.0	223
square-ctr	5	2	18.4	171	5	1209.4	2111	5	17.0	258
logistics	3	3	0.0	24	1	7.5	160	3	0.1	41
look-and-grab	18	11	0.4	31	7	5.2	9	15	0.1	11
push-to	9	6	180.2	227	7	67.8	89	8	65.1	219
Raos-keys	3	2	0.0	16	2	0.5	22	1	0.7	21
ring	3	3	0.1	55	1	1546.3	39	3	0.6	41
Uts-k	6	6	4.8	82	6	9.6	92	4	13.1	100

T1 compared with conf. planners *T0* and *DNF* (To, Son & Pontelli, 2010)

T1 has a **higher coverage** than both *T0* and *DNF*.

Heuristics comparison: $h_C(b)$ vs. $h_K(b)$

Domain	I	h_C				h_K				T1			
		S	T	E	L	S	T	E	L	S	T	E	L
Bomb	9	7	71	4k	101	7	11	773	101	8	2	100	101
Cube(Ctr)	12	6	84	32k	188	10	1	890	61	12	0.1	61	58
Cube(Cor)	11	8	92	219k	271	10	4	26k	88	11	12	15k	269
Dispose	11	7	664	8k	349	9	57	2k	190	8	134	1k	491
Logistics	4	2	0.2	546	30	2	544	1613k	30	4	0.1	554	78
Ring	7	6	1	1k	17	5	571	58k	17	8	0.2	214	31
UTS-k	15	15	0.06	26	7	2	0.05	154	7	13	0.04	10	9

- The **combination** of h_C and h_K in T1 performs generally better

- 1 A method to define and compute a **sample set** $S \subseteq S_o$ s.t.
 - $|S|$ is exponential in i
 - translation $K_S^i(P)$ based on S is sound and complete if $width(P) \leq i$
- 2 A **Planner** $T1$ based on new translation, extended with certainty heuristic, **competitive** with state-of-the-art.

All good things come to an End

Thank you!