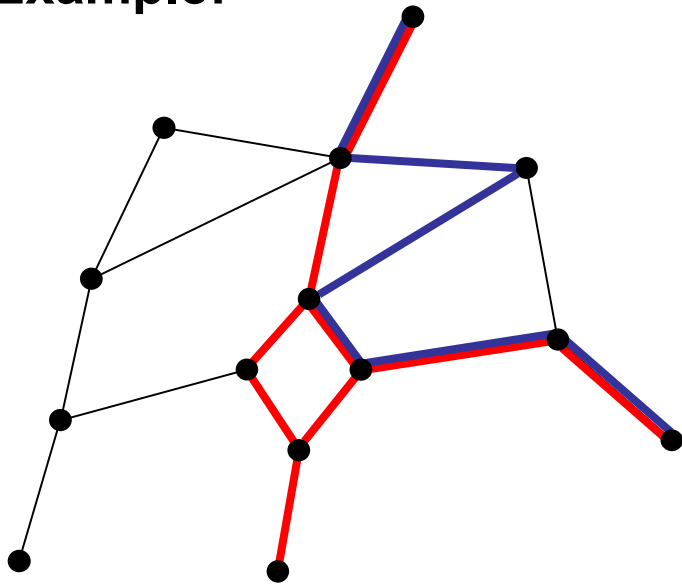


Closed Pattern Mining in Strongly Accessible Set Systems

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given finite set E ,
membership oracle $M_{\mathcal{F}}: 2^E \rightarrow \{0,1\}$ (defining $\mathcal{F} \subseteq 2^E$),
closure operator $\rho: \mathcal{F} \rightarrow \mathcal{F}$,
list the family of ρ -closed sets $\rho(\mathcal{F}) = \{C \in \mathcal{F}: C = \rho(C)\}$

Example:



street network: graph G

transaction database \mathcal{D} : subsets of $E(G)$

- E : edge set $E(G)$
- \mathcal{F} : $\{X \subseteq E(G) : X \text{ is freq. and } \textit{connected}\}$
- $\rho: X \mapsto$ connected component of X in

$$\bigcap \{T \in \mathcal{D}: T \supseteq X\}$$

main result:

problem can be solved efficiently with a simple DFS algorithm for a general class of set systems (which for instance do not satisfy anti-monotonicity)

set system (E, \mathcal{F}) is **strongly accessible** if $\emptyset \in \mathcal{F}$ and for every $X, Y \in \mathcal{F}$ satisfying $X \subset Y$, there exists an $e \in Y \setminus X$ such that $Y \setminus \{e\} \in \mathcal{F}$.

\Rightarrow there is a sequence $X = X_0, X_1, \dots, X_k = Y$ s.t. $|X_i \setminus X_{i-1}| = 1$ for $i = 1, \dots, k$

formally the following statement holds:

Thm: For any finite strongly accessible set system (E, \mathcal{F})

(i) given by a membership oracle M computable in time T_M ,

(ii) for any closure operator $\rho: \mathcal{F} \rightarrow \mathcal{F}$, computable in time T_ρ

the family of closed sets can be listed with delay $(|E|^2 T_M T_\rho)$.

Details at the poster!